

Lecture 1.2, MATH-42021/52021 Graph Theory and Combinatorics.

Artem Zvavitch

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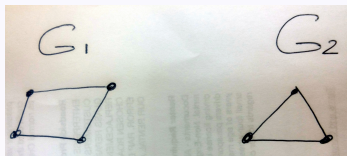
June, 2016.

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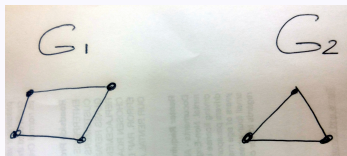
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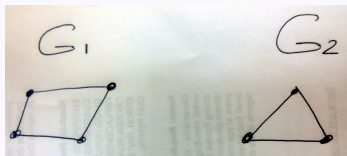


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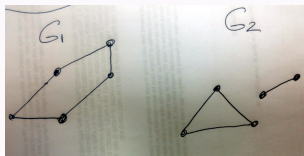


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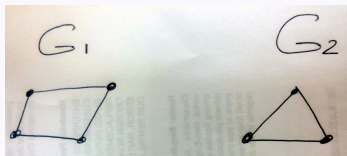


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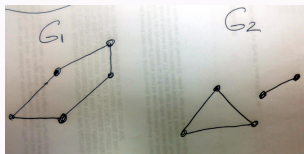


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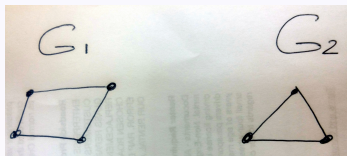


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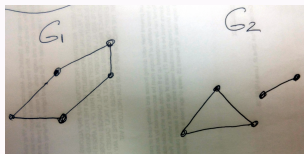


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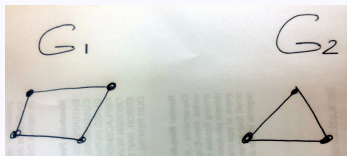


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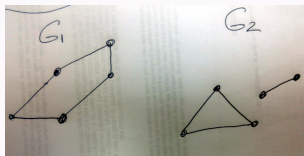


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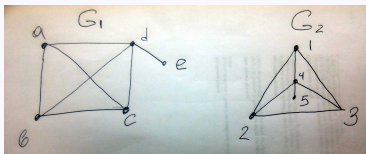


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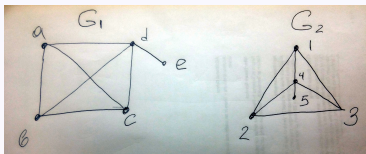
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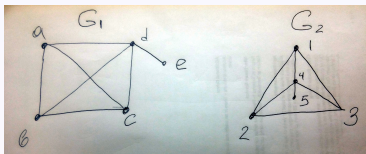


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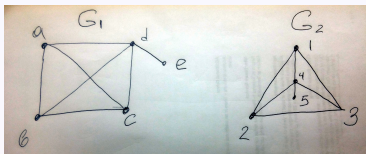


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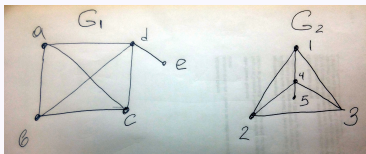


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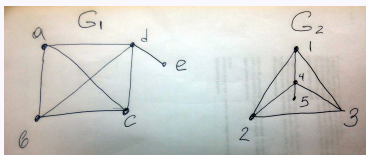
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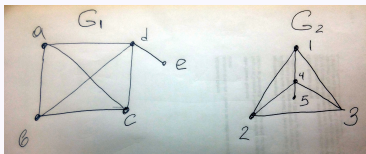
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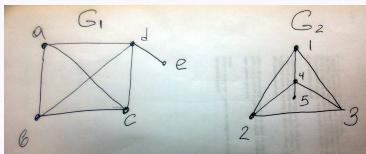
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Next it is not hard to check that f is isomorphism, indeed this is one to one map and we can check the property of preserving the adjacency (for example f maps (e, d) to $(5, 4)$).

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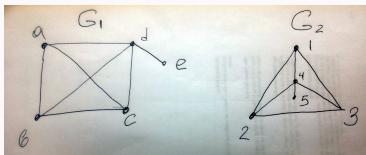
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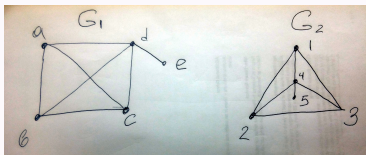


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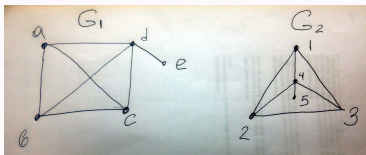
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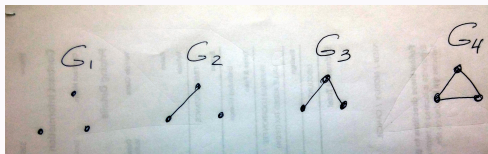
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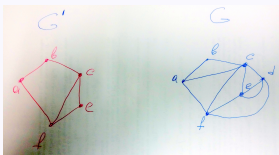
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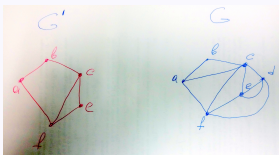
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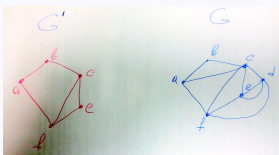
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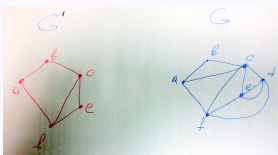
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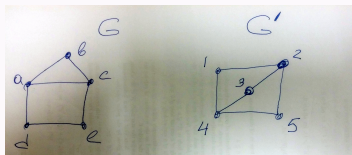
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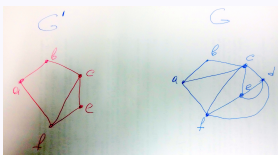


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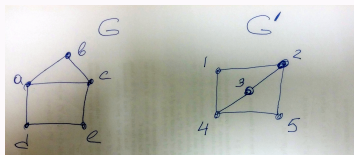


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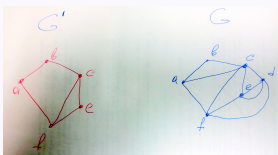
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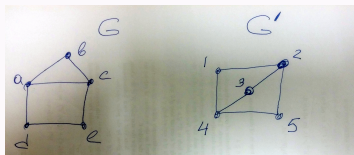
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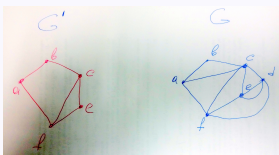


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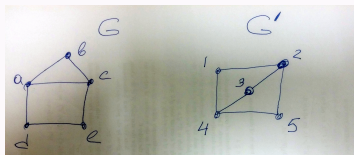


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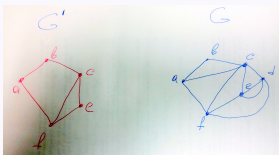
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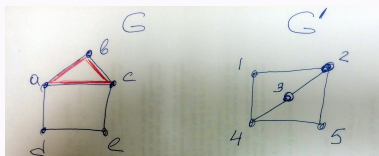
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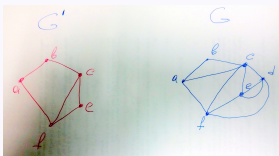


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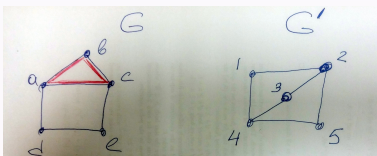
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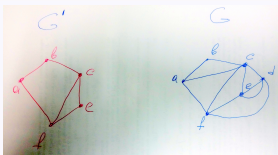


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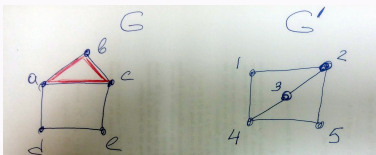
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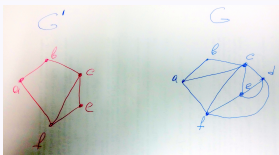


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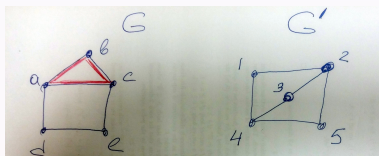
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A subgraph is a graph formed by a subset of vertices and edges of a larger graph.



Subgraphs are super useful when we study isomorphisms. Remember, the isomorphisms must preserve "structure" thus must preserve subgraphs. So do you think those two graphs are isomorphic:



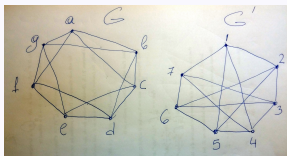
- Vertex count - the same.
- Edge count - the same.
- Degree count is the same. Both graphs 3 vertices of degree 2 and 2 vertices of degree 3.

What about now? You may see that G has a subgraph of 3 vertices $\{a, b, c\}$ which is a circuit, there is no circuit of length 3 in G' . So the answer is NO those graphs are not isomorphic.

NOTICE: this is an example of two graphs, with the same, vertex, edge and degree count which are NOT isomorphic.

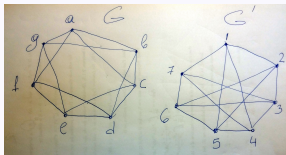
Another example

Are those two graphs isomorphic?



Another example

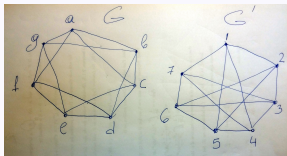
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- Vertex count - the same.

Another example

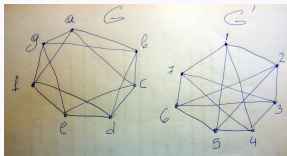
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- Vertex count - the same.
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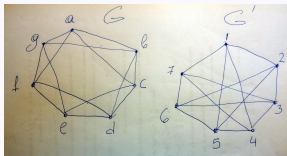
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- Vertex count - the same.
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- More over, degree count is the same.

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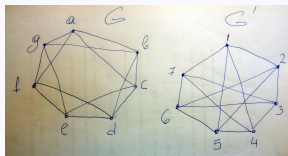
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- More over, degree count is the same.
- Trying to look for some contradiction in subgraphs no luck.

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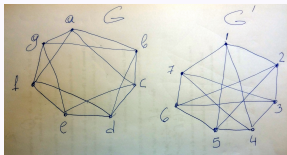


- Vertex count - the same.
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- More over, degree count is the same.
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So we need just to start working out isomorphism and see how it goes.....

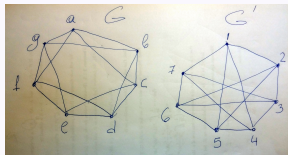
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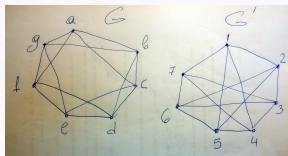
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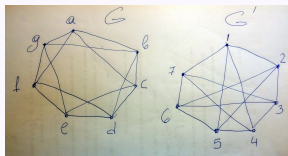
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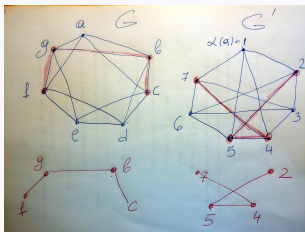
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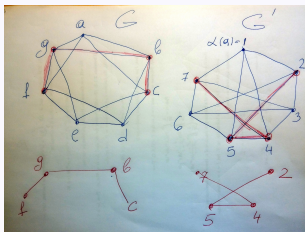
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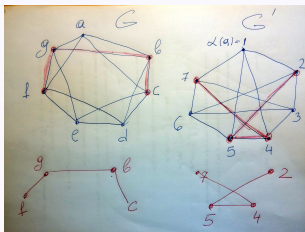
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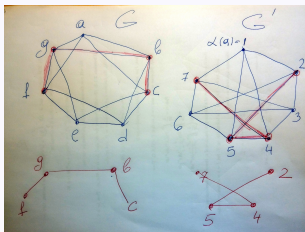
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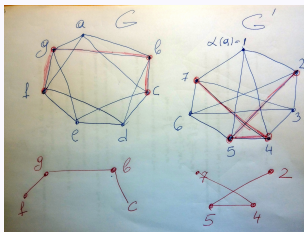
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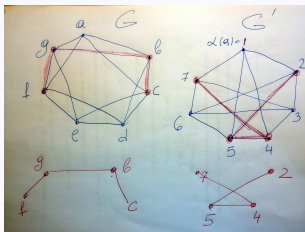
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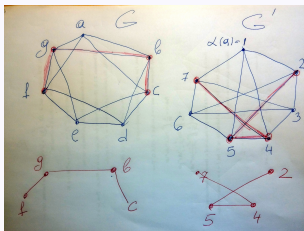
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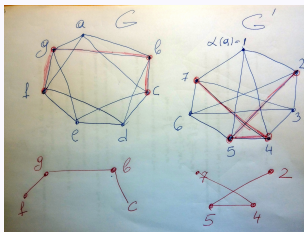
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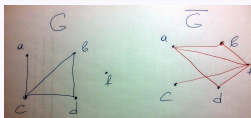
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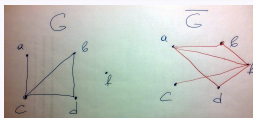
Yet, another example + a cool idea

Consider a graph $G = (V, E)$, its **complement** is a graph $\bar{G} = (V, E')$ with the same set of vertices but now with edges between exactly those pairs of vertices that were not adjacent in G .



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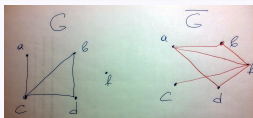
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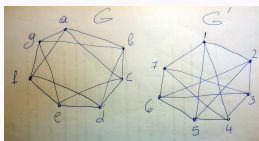
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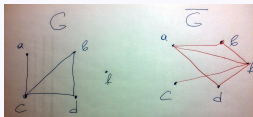


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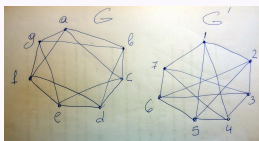


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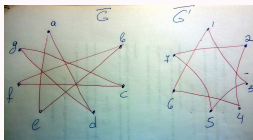
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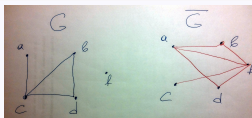


Now lets draw a complement graphs

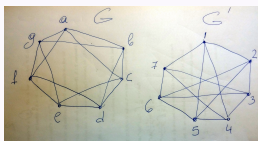


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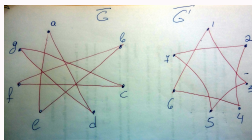
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Now lets draw a complement graphs



Both of them are circuits of length 7, thus isomorphic, and solution became almost trivial with this cool trick!!!